Model-Checking Based Verification of Cyber-Physical Systems with Alternating Signal Temporal Logic

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Motivation

Cyber-physical systems (CPS) utilize software while at the same time being able to sense or interact with the real world. We could model the system "Environment and CPS" and then verify the CPS by model checking a specification against the system.

Bounded Model Checking

Extension of Pre-Image based algorithm for ATL

Algorithm 1 ASTL symbolic model-checking **Input:** timed game \mathcal{T} , ASTL formula φ **Output:** boolean *true* or *false*

Signal Temporal Logic

 $\varphi := s \sim c \mid \top \mid \neg \varphi \mid (\varphi \lor \varphi) \mid (\varphi \mathbf{U}_I \varphi)$

Real-time • Real-valued • Linear-time

Alternating-Time Temporal Logic

 $\varphi := p \mid \top \mid \neg \varphi \mid (\varphi \lor \varphi) \mid \langle \! \langle A \rangle \! \rangle \bigcirc \varphi \mid \langle \! \langle A \rangle \! \rangle \square \varphi \mid \langle \! \langle A \rangle \! \rangle (\varphi \mathbf{U} \varphi)$

Discrete-time • Discrete-valued • Branching-time Can express existence of strategies for players

Alternating Signal Temporal Logic

$\varphi := s \sim c \mid \top \mid \neg \varphi \mid (\varphi \lor \varphi) \mid (\varphi \mathbf{U}_I \varphi)$

Real-time • Real-valued • Branching-time

for φ' in $\operatorname{Sub}(\varphi)$ do	
$\mathbf{case} \varphi' = \top$	
$[\varphi'] \leftarrow Q_I$	
case $\varphi' = x \sim c$	Rules
$[\varphi'] \leftarrow \operatorname{Reg}_{\mathcal{T}}(x \sim c)$	1. System "timed game" as game-
case $\varphi' = \neg \theta$	structure extension of timed automata
$[\varphi'] \leftarrow Q_T \setminus [\theta]$	2. Only non-zeno structures
$[\mathcal{Y}] \land \mathcal{Q}I \land [\mathcal{V}]$	3. No infinite intervals in specification
$\mathbf{case} \ \varphi' = (\theta_1 \lor \theta_2)$	
$[\varphi'] \leftarrow [\theta_1] \cup [\theta_2]$	
case $\varphi' = \langle\!\langle A \rangle\!\rangle \left(\theta_1 \mathbf{U}_I \theta_2\right)$	
$[\varphi'] \leftarrow \operatorname{Pre}^*_{\mathcal{T},I}(A, [\theta_2], [$	$ heta_1])$
end for	
$\mathbf{return} \ q_0 \in [\varphi]$	

Proof Ideas

Correctness: Induction over ASTL semantics **Computability:**

Can express existence of strategies for players

Timed Games

Timed automaton extension with multiple players and actions for said players activation transitions



- 1. Split pre-image computation into series of discrete (action) and continuous (time) steps
- 2. Show that each step is computable, and a preimage computation involves finitely many steps
- 3. Show that any pre-image can be symbolically represented as regions in state-space
- 4. Final inclusion check of algorithm through translation of symbolic representation to firstorder logic formulas over the theory of realclosed fields, quantifier elimination computable (Tarski-Seidenberg Theorem)

Caveats and Further Work

- Infeasible runtime (quantifier-elimination in the theory of real-closed fields) – really necessary?
- ASTL* in relation to ASTL (analogously to ATL*)
- Robust semantics for ASTL

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